

Comparison of Routing Metrics for Static Multi-Hop Wireless Networks

Richard Draves, Jitendra Padhye and Brian Zill

Microsoft Research

Multi-hop Wireless Networks

	Static	Mobile
Motivating scenario	Community wireless networks (“Mesh Networks”)	Battlefield networks
Key challenge	Improving network capacity	Handling mobility, node failures, limited power.

Routing in Multi-hop Wireless Networks

- Mobile networks:
 - Minimum-hop routing (“shortest path”)
 - DSR, AODV, TORA

- Static networks:
 - Minimum-hop routing tends to choose long, lossy wireless links
 - Taking more hops on better-quality links can improve throughput

[De Couto et. al., HOTNETS 2003]

Link-quality Based Routing

- Metrics to measure wireless link quality:
 - Signal-to-Noise ratio
 - Packet loss rate
 - Round trip time
 - Bandwidth
 - ...

Our paper: experimental comparison of performance of three metrics in a 23 node, indoor testbed.

Contributions of our paper

- Design and implementation of a routing protocol that incorporates notion of link quality
 - Link Quality Source Routing (LQSR)
 - Operates at layer “2.5”
- Detailed, “side-by-side” experimental comparison of three link quality metrics:
 - Per-hop Round Trip Time (RTT) [*Adya et al 2004*]
 - Per-hop Packet Pair (PktPair)
 - Expected Transmissions (ETX) [*De Couto et al 2003*]

Summary of Results

- ETX provides best performance
- Performance of RTT and PktPair suffers due to self-interference
- PktPair suffers from self-interference only on multi-hop paths

Outline of the rest of the talk

- LQSR architecture (brief)
- Description of three link quality metrics
- Experimental results
- Conclusion

LQSR Architecture

- Source-routed, link-state protocol
 - Derived from DSR
- Each node measures the quality of links to its neighbors
- This information propagates throughout the mesh
- Source selects route with best cumulative metric
- Packets are source-routed using this route

Link Quality Metrics

- Per-hop Round Trip Time (**RTT**)
- Per-hop Packet-Pair (**PktPair**)
- Expected transmissions (**ETX**)
- Minimum-hop routing (**HOP**)
 - Binary link quality

Metric 1: Per-hop RTT

- Node periodically pings each of its neighbors
 - Unicast probe/probe-reply pair
- RTT samples are averaged using TCP-like low-pass filter
- Path with least sum of RTTs is selected

Metric 1: Per-hop RTT

- Advantages
 - Easy to implement
 - Accounts for link load and bandwidth
 - Also accounts for link loss rate
 - 802.11 retransmits lost packets up to 7 times
 - Lossy links will have higher RTT
- *Disadvantages*
 - Expensive
 - **Self-interference due to queuing**

Metric 2: Per-hop Packet-Pair

- Node periodically sends two back-to-back probes to each neighbor
 - First probe is small, second is large
- Neighbor measures delay between the arrival of the two probes; reports back to the sender
- Sender averages delay samples using low-pass filter
- Path with least sum of delays is selected

Metric 2: Per-hop Packet-Pair

- *Advantages*
 - Self-interference due to queuing is not a problem
 - Implicitly takes load, bandwidth and loss rate into account
- *Disadvantages*
 - More expensive than RTT

Metric 3: Expected Transmissions

- **Estimate number of times a packet has to be retransmitted on each hop**
- Each node periodically broadcasts a probe
 - 802.11 does not retransmit broadcast packets
- Probe carries information about probes received from neighbors
- Node can calculate loss rate on forward (P_f) and reverse (P_r) link to each neighbor

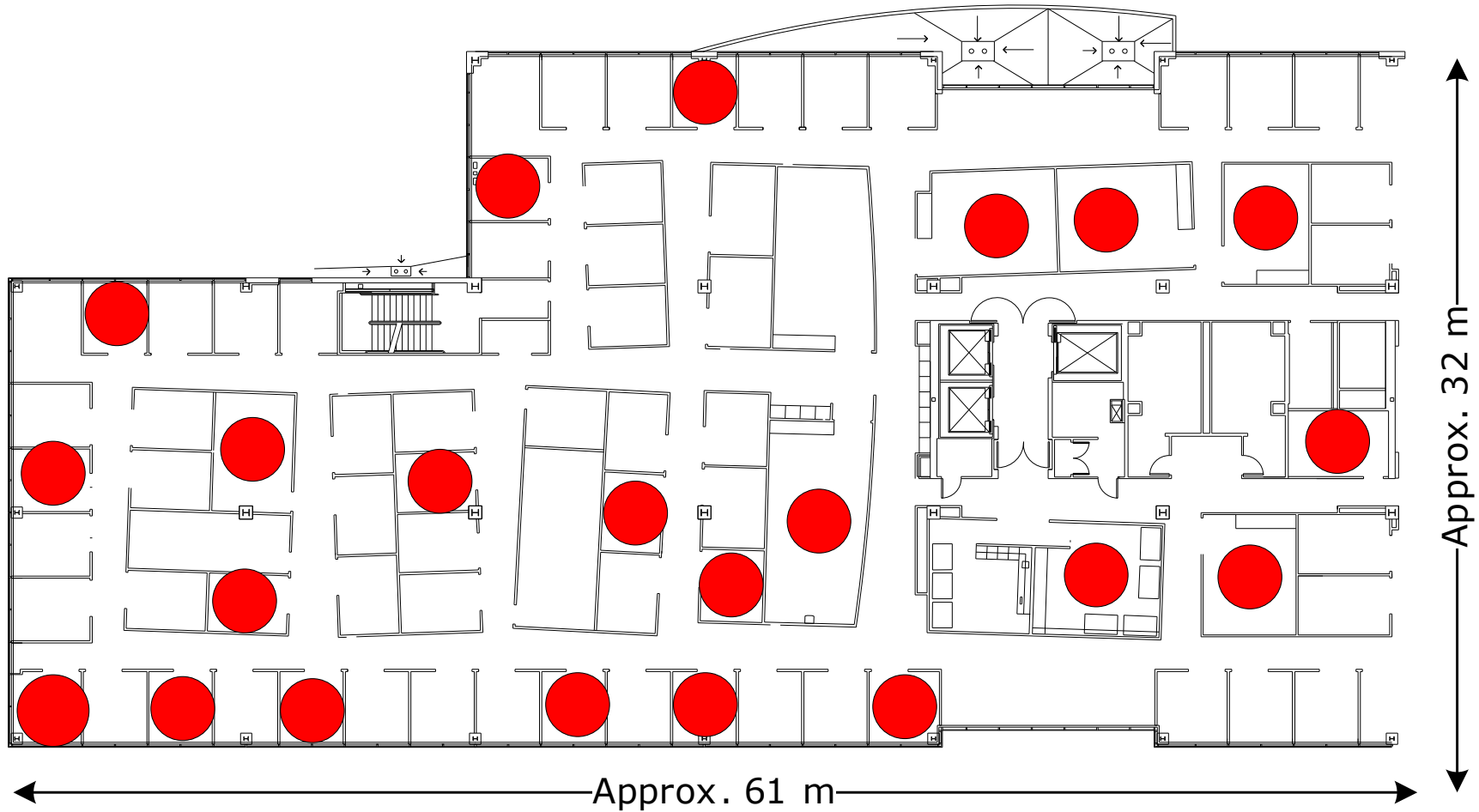
$$ETX = \frac{1}{(1 - P_f) * (1 - P_r)}$$

- **Select the path with least total ETX**

Metric 3: Expected Transmissions

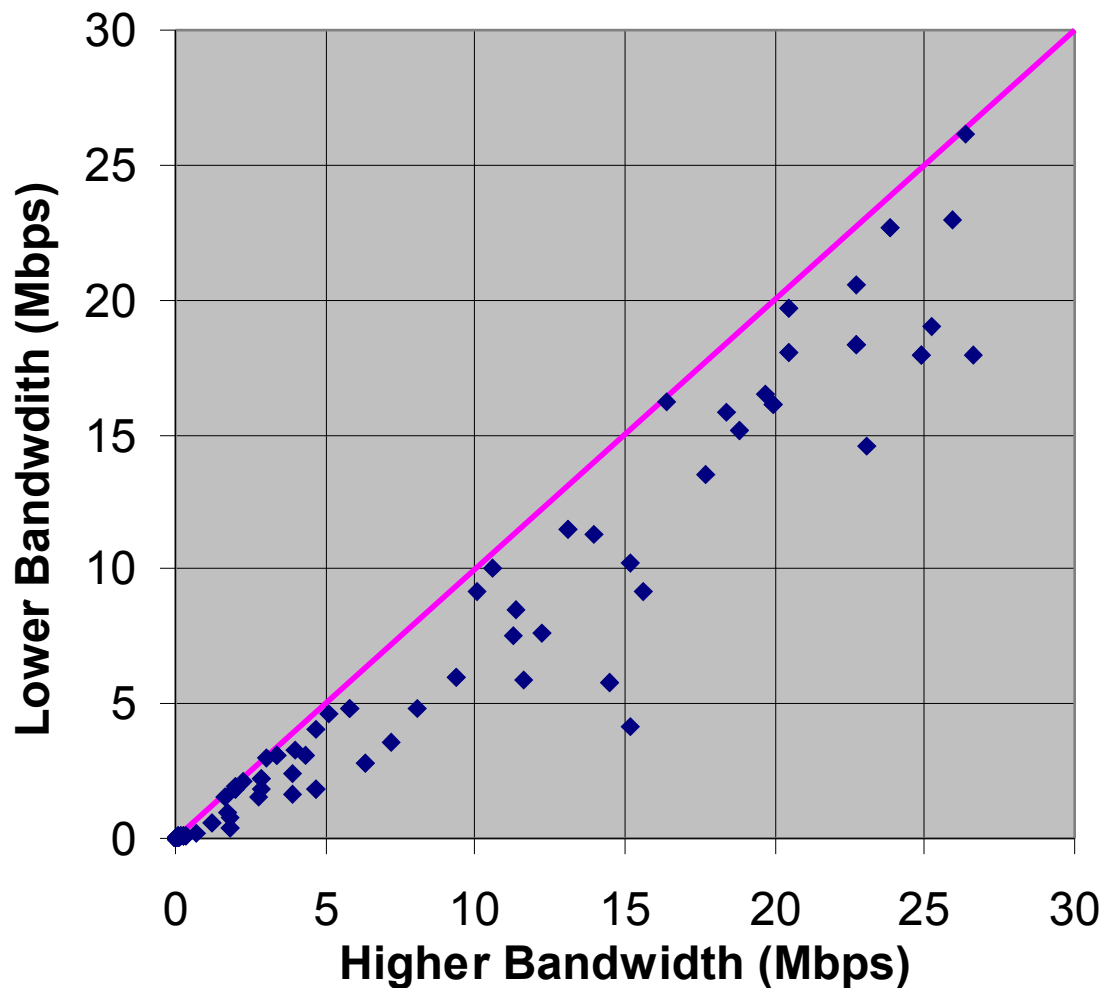
- *Advantages*
 - Low overhead
 - Explicitly takes loss rate into account
- *Disadvantages*
 - Loss rate of broadcast probe packets is not the same as loss rate of data packets
 - Probe packets are smaller than data packets
 - Broadcast packets are sent at lower data rate
 - Does not take data rate or link load into account

Mesh Testbed



**23 Laptops running Windows XP.
802.11a cards: mix of Proxim and Netgear.
Diameter: 6-7 hops.**

Link bandwidths in the testbed



- Cards use Autorate

- Total node pairs:
 $23 \times 22 / 2 = 253$

- 90 pairs have non-zero bandwidth in both directions.

Bandwidths vary significantly; lot of asymmetry.

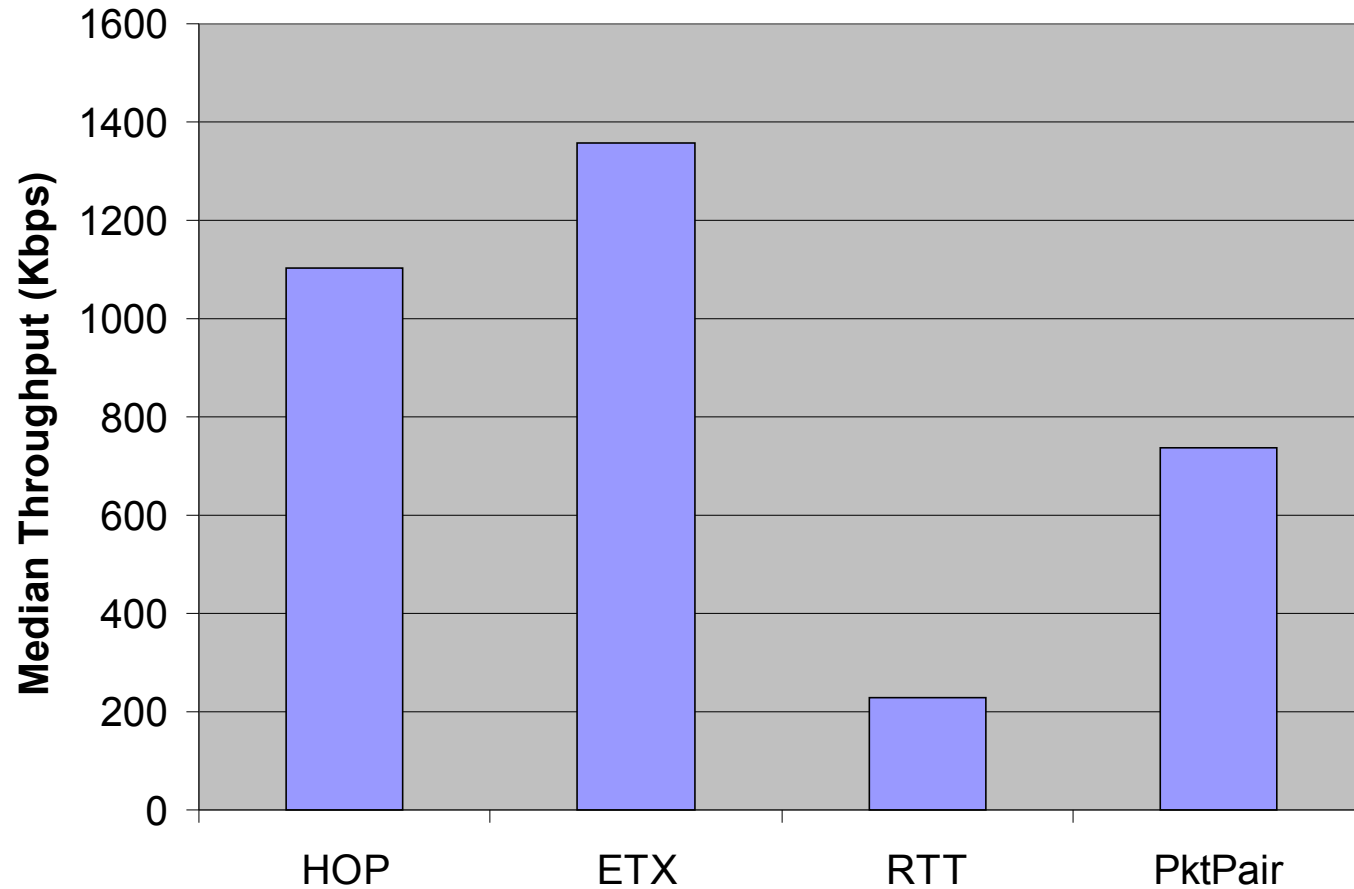
Experiments

1. Bulk-transfer TCP Flows
-
-
4. Impact of mobility

Experiment 1

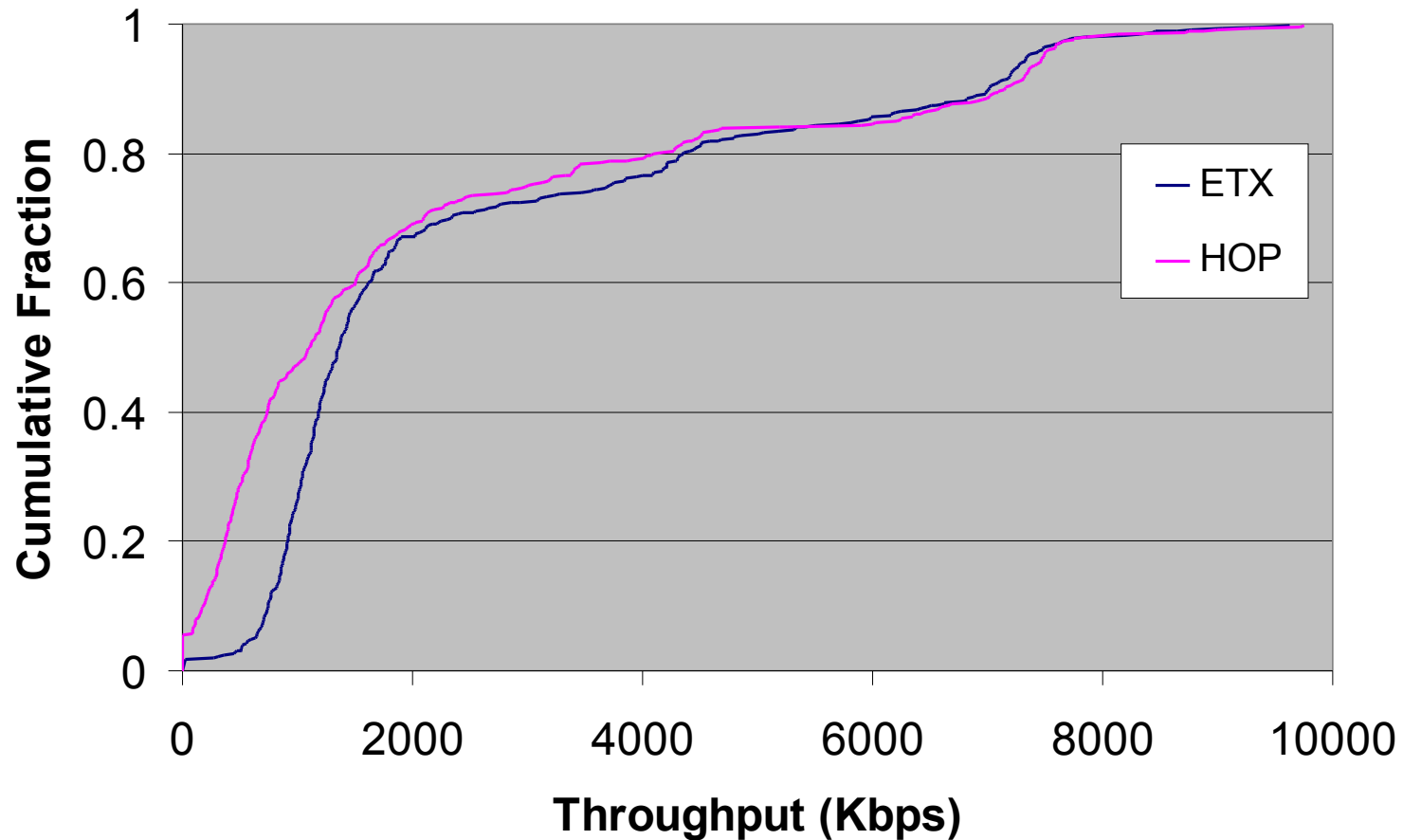
- 3-Minute TCP transfer between each node pair
 - $23 \times 22 = 506$ pairs
 - 1 transfer at a time
 - Long transfers essential for consistent results
- For each transfer, record:
 - Throughput
 - Number of paths
 - Path may change during transfer
 - Average path length
 - Weighted by fraction of packets along each path

Median Throughput



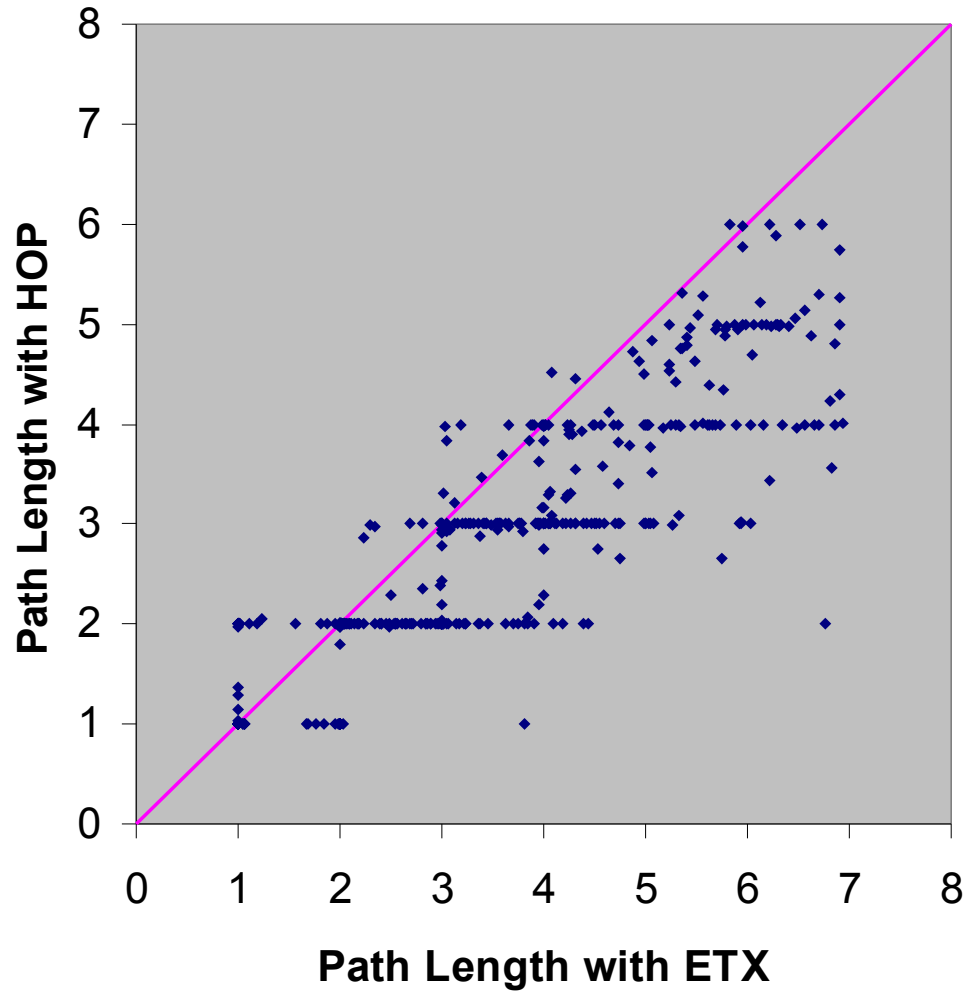
ETX performs best. RTT performs worst.

Why does ETX perform well?



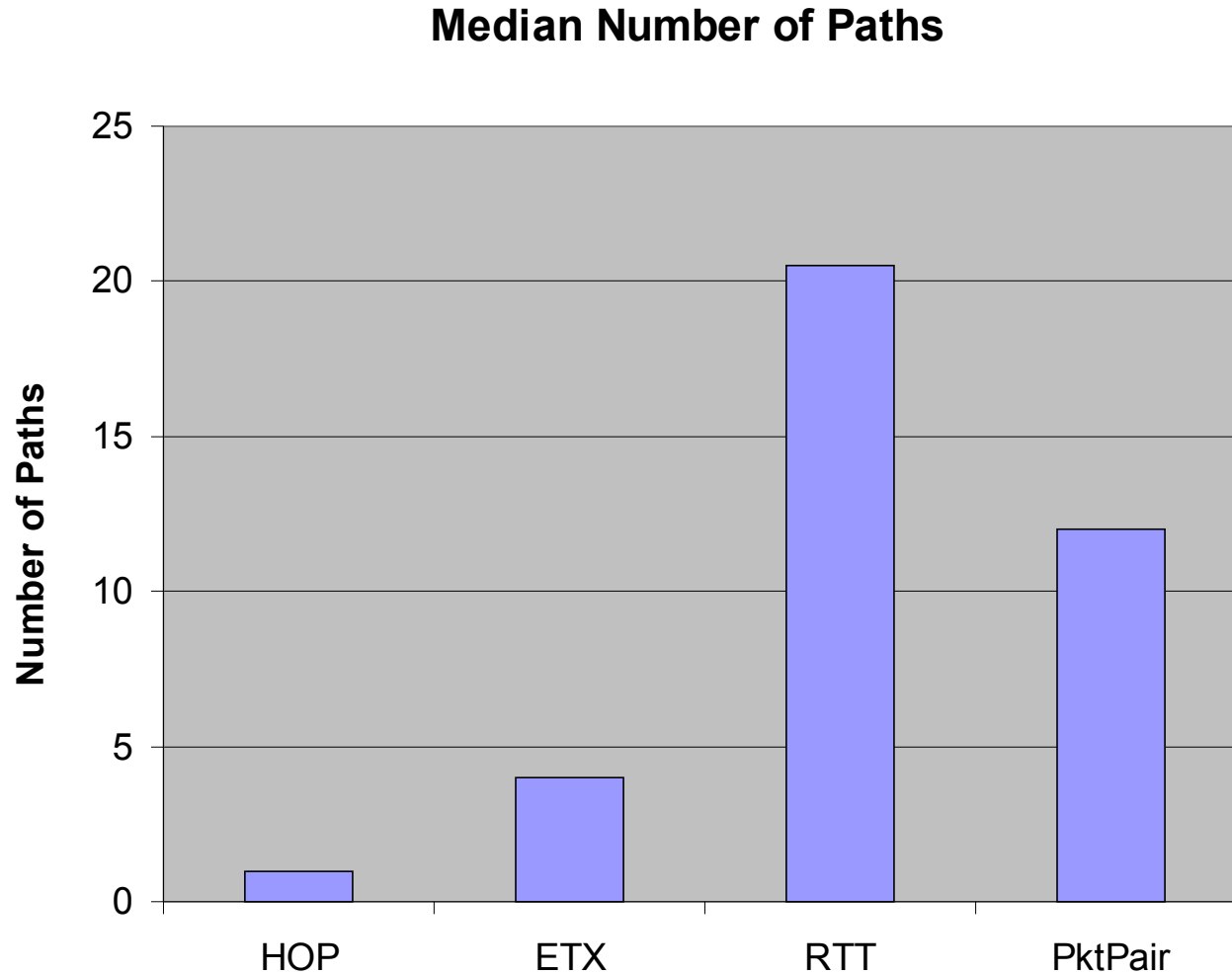
ETX performs better by avoiding low-throughput paths.

Impact on Path Lengths



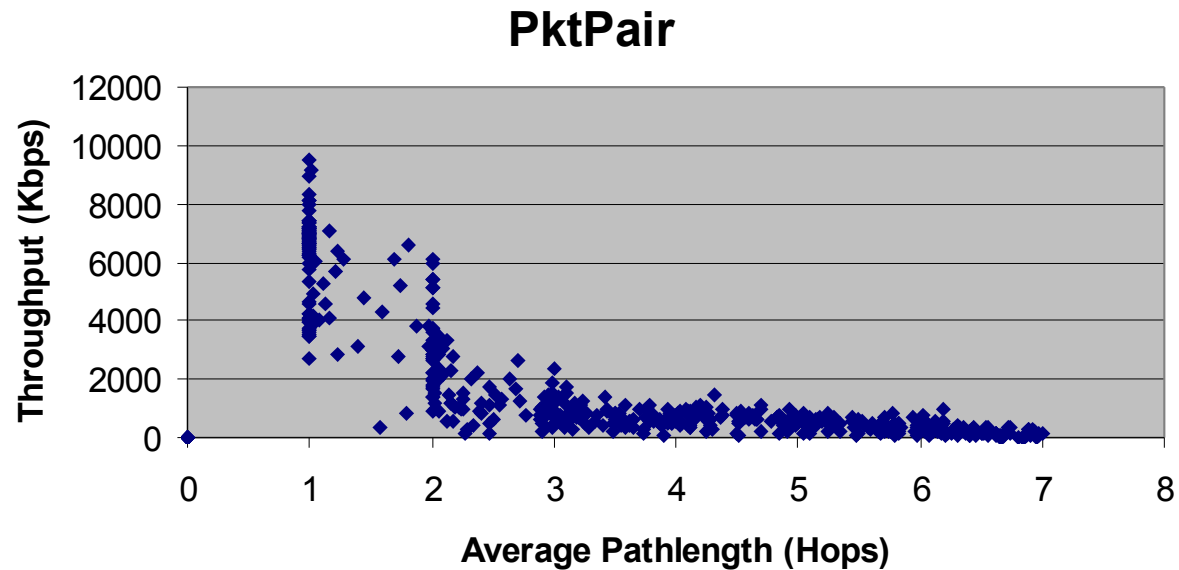
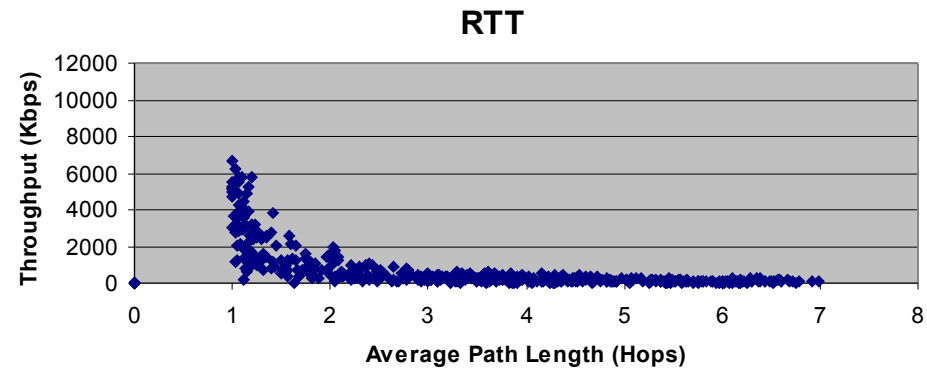
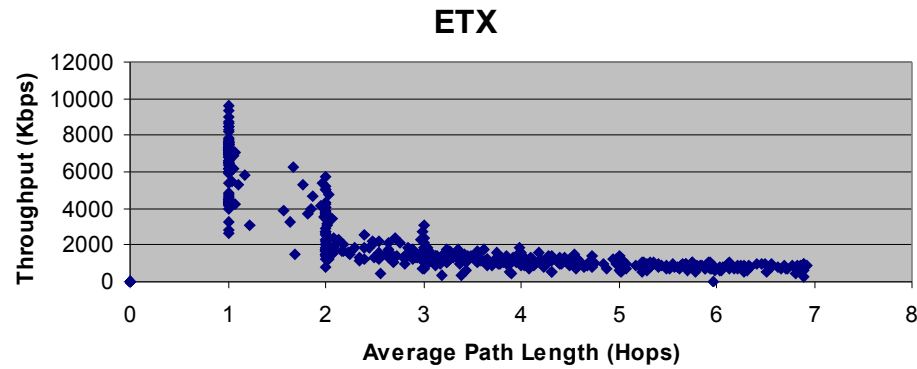
Path length is generally higher under ETX.

Why does RTT perform so poorly?



RTT suffers heavily from self-interference

What ails PktPair?



PktPair suffers from self-interference only on multi-hop paths.

Summary

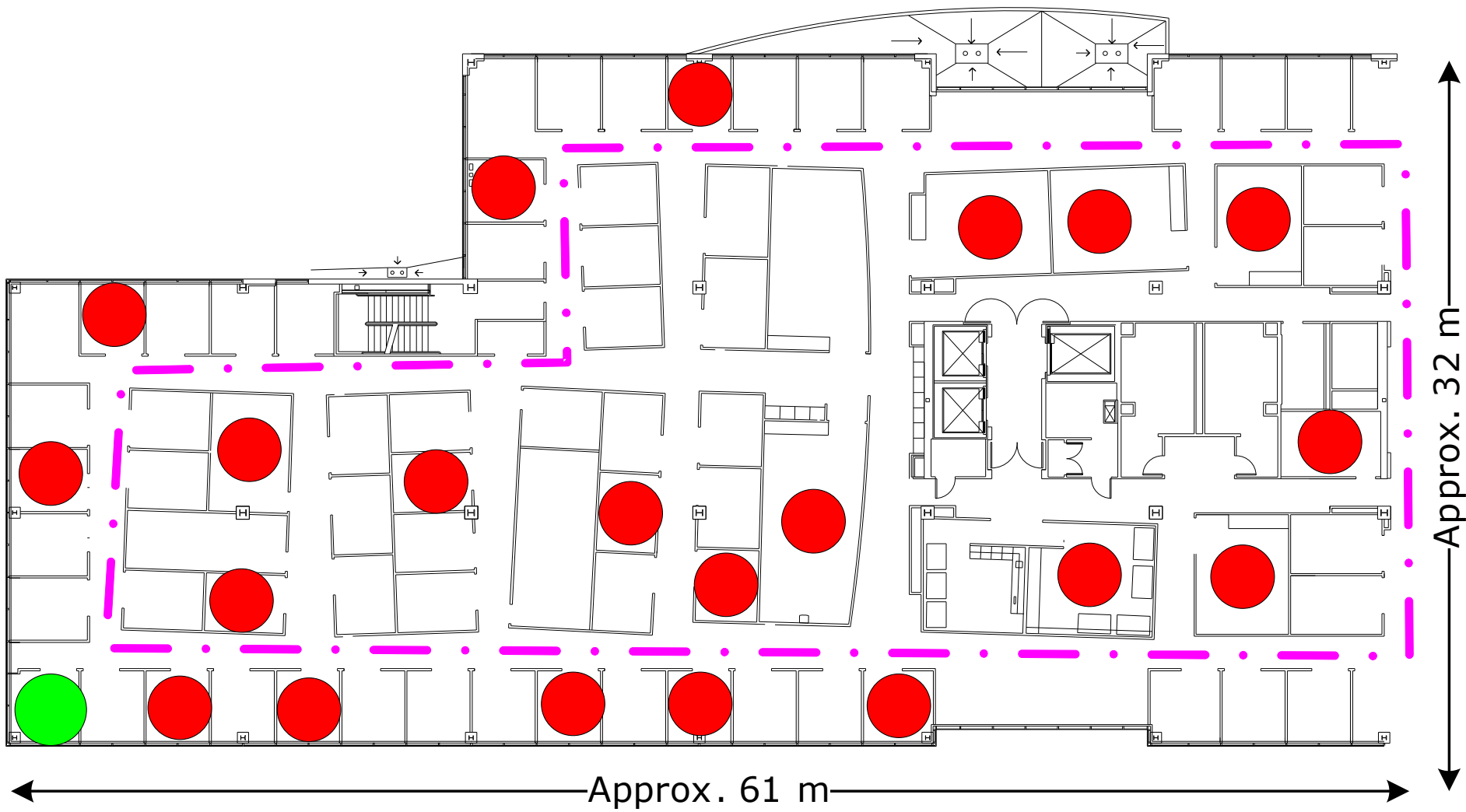
- ETX performs well despite ignoring link bandwidth
- Self-interference is the main reason behind poor performance of RTT and PktPair.

Similar results for multiple simultaneous flows.

Experiment 2

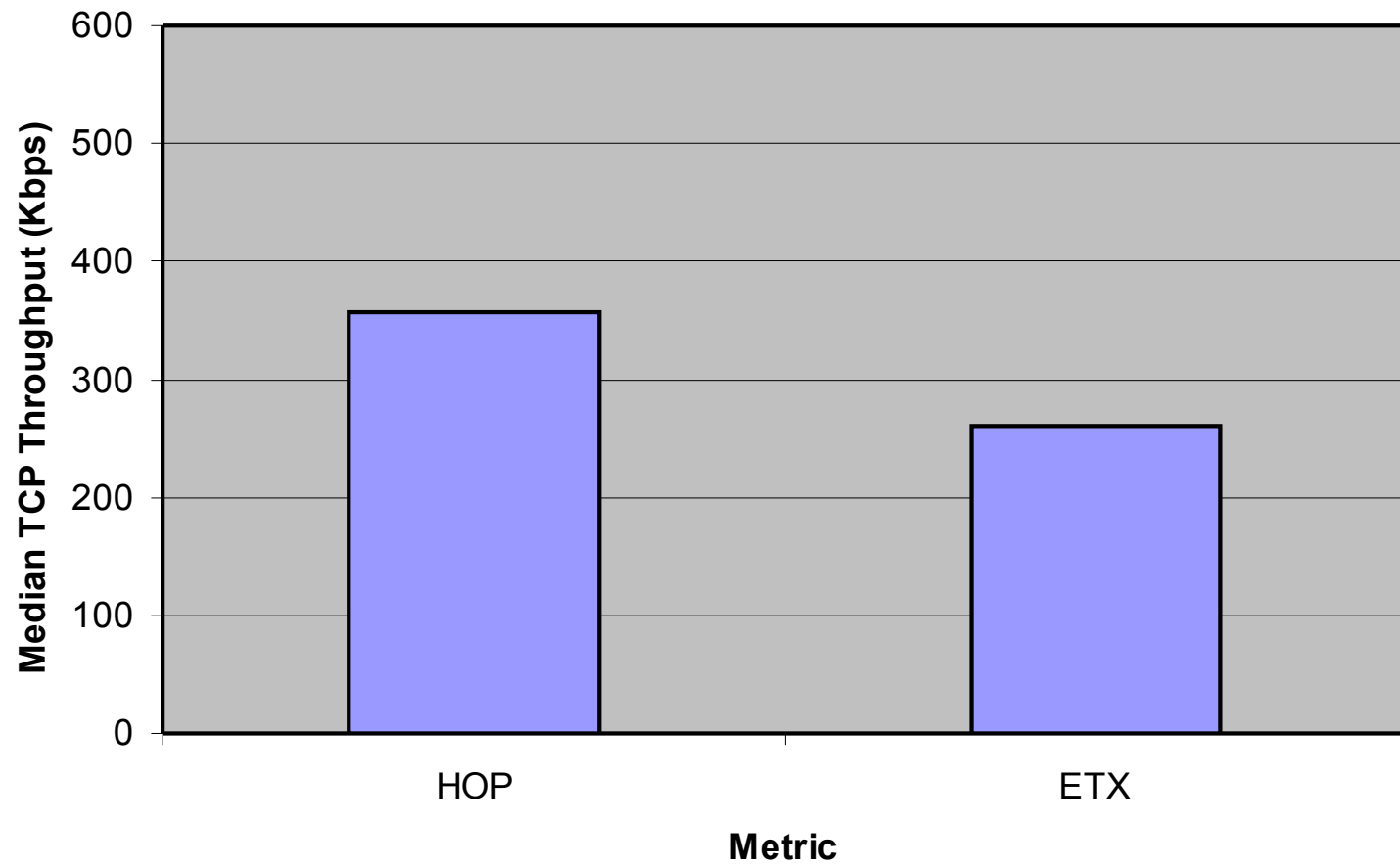
- Walk slowly around network periphery for 15 minutes with a laptop
- Mobile laptop is the sender, a corner node is receiver
- Repeated 1-minute TCP transfers

Testbed Layout



Approx. 61 m

Approx. 32 m



Shortest path routing is best in mobile scenarios?

Conclusions

- ETX metric performs best in static scenarios
- RTT performs worst
- PacketPair suffers from self-interference on multi-hop paths
- Shortest path routing seems to perform best in mobile scenarios
 - Metric-based routing does not converge quickly?

Ongoing/Future work

- Explicitly take link bandwidth into account
- Support for multiple heterogeneous radios per node
 - To appear in MOBICOM 2004
- Detailed study of TCP performance in multi-hop networks
- Repeat study in other testbeds

For more information

<http://research.microsoft.com/mesh/>

Source code, binaries, tech reports, ...

Backup slides

LQSR Architecture

- Implemented in a shim layer between Layer 2 and 3.
- The shim layer acts as a virtual Ethernet adapter
 - Virtual Ethernet addresses
 - Multiplexes heterogeneous physical links
- Advantages:
 - Supports multiple link technologies
 - Supports IPv4, IPv6 etc *unmodified*
 - Preserves the link abstraction
 - Can support any routing protocol

- Architecture:

IPv4	IPv6	IPX
Mesh connectivity Layer with LQSR		
Ethernet	802.11	802.16

- Header Format:

Ethernet	MCL	Payload: TCP/IP, ARP, IPv6...
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Web transfers

- Simulated Web transfer using Surge
- One node serves as web server
- Six nodes along periphery act as clients
- Results: ETX reduces latency by 20% for hosts that are more than one hop away from server.

Static Multi-hop Wireless Networks

- Motivating scenario:
 - Community wireless networks (“Mesh Networks”)
 - Very little node mobility
 - Energy not a concern
- Main Challenge:
 - Improve Network capacity
- Minimum-hop count routing is inadequate
 - Tends to choose long, lossy wireless links [*De Couto et. al., HOTNETS 2003*]

“Traditional” Multi-hop Wireless Networks

- Envisioned for mobility-intensive scenarios
- Main concerns:
 - Reduce Power consumption
 - Robustness in presence of mobility, link failures
- Routing:
 - Minimum-hop routing (“shortest path”) with various modifications to address power and mobility concerns
 - DSR, AODV, TORA