

Mixing Email with BABEL'96

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Introduction

- The most critical network building block for anonymity: Mix by Chaum'81
- Absolute anonymity: dining cryptographers algorithm and protocol by Chaum'88
 - Impractical due to the large protocol overhead
 - Hard to provide secure pairwise channels and a synchronized broadcast channel
- Advancement being sought

Related work

- Anonymous Email systems
 - Penet: Finland 1990s
 - Cypherpunk: 1990s – now?
 - Mixmaster: 1990s – now?
 - Implementations of the concept of Mix by a single individual or a group of contributors
- Systematic improvements needed

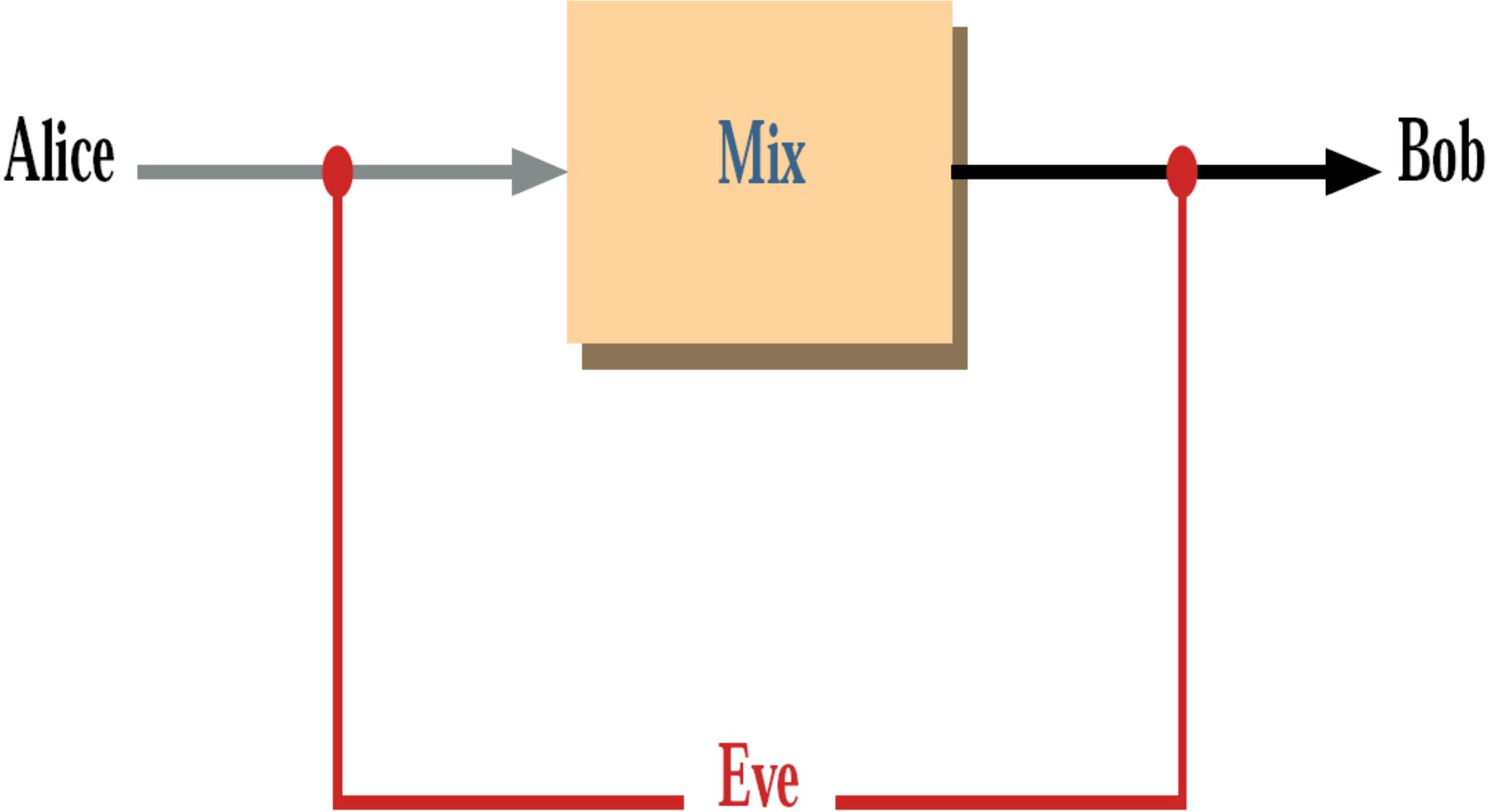
Desired properties (requirements)

- Email systems accommodate anonymity
- Hard to determine the sender
- Recipient can reply the anonymous sender
- End-point anonymity preserved, intermediate mixes are not maximally trusted
- Infrastructure resistant to active attacks
- Sender gets an anonymous confirmation
- Low noise

Notation

M	message; sequence of ASCII bits
$E_x(M)$	encryption of M with X 's public key
$D_x(M)$	decryption of M with X 's private key
$K\{M\}$	conventional encryption of M with key K
(M_1, M_2)	concatenation of M_1 and M_2
A_x	X 's email address.
$[M]^\Omega$	padding string M to length Ω (by appending random bits)
$\lfloor M \rfloor^\Omega$	trimming string M to length Ω (by removing trailing bits)

Mix, revisted



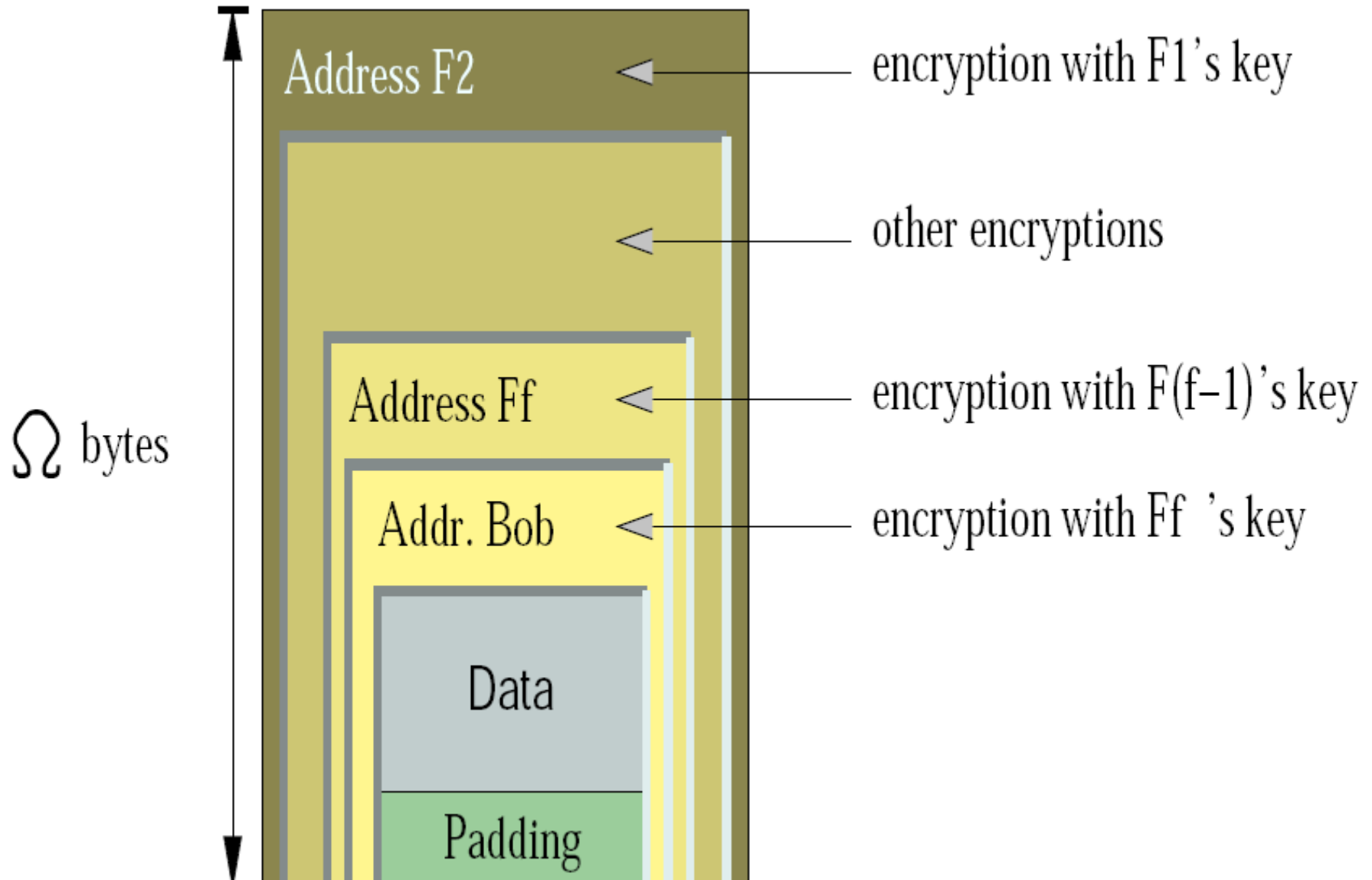
Attacks

- Passive
 - Content correlation:
 - uniform length, padding, nonce
 - Time correlation:
 - Regular vs. interval batching
- Active
 - Isolate and Identify
 - Message replay:
 - time stamp, message identifier
 - Cascading or chaining mixes

BABEL

- Forward path
 - Composition by sender
 - Processing by mixes
 - What does a mix know?

Forward message



Forward message

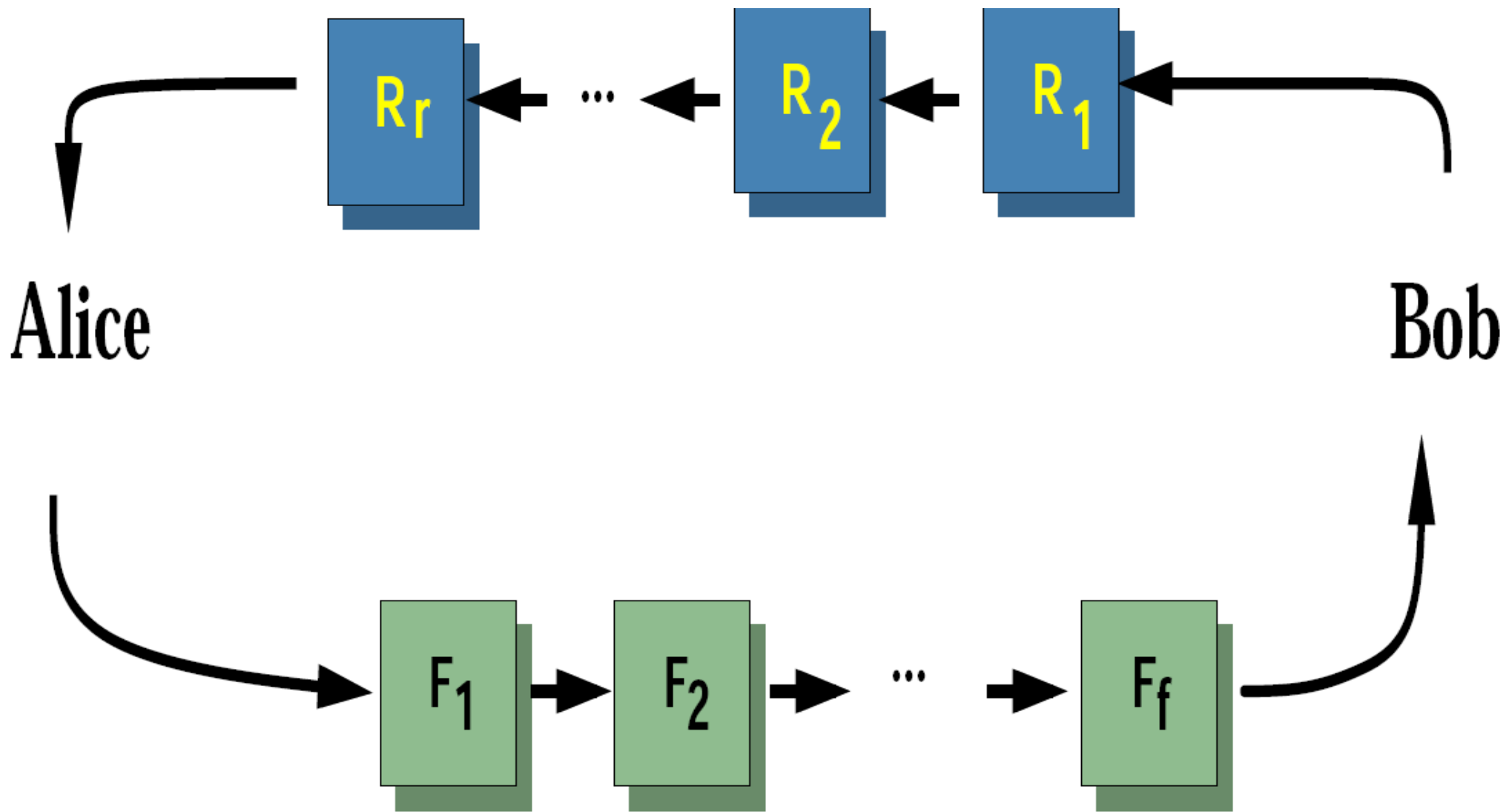
$x_f =$

$$E_{F_1}(\mathcal{A}_{F_2}, E_{F_2}(\dots E_{F_f-1}(\mathcal{A}_{F_f}, E_{F_f}(\mathcal{A}_{\text{Bob}}, [M]^\Omega)) \dots))$$

BABEL

- Return path
 - Creating RPI
 - Replying by recipient
 - Reply processing by BABEL
 - Handling replies at the originator
 - Two-way anonymous conversation
 - Security of replies
 - Inter-mix Detours
 - Indirect replies

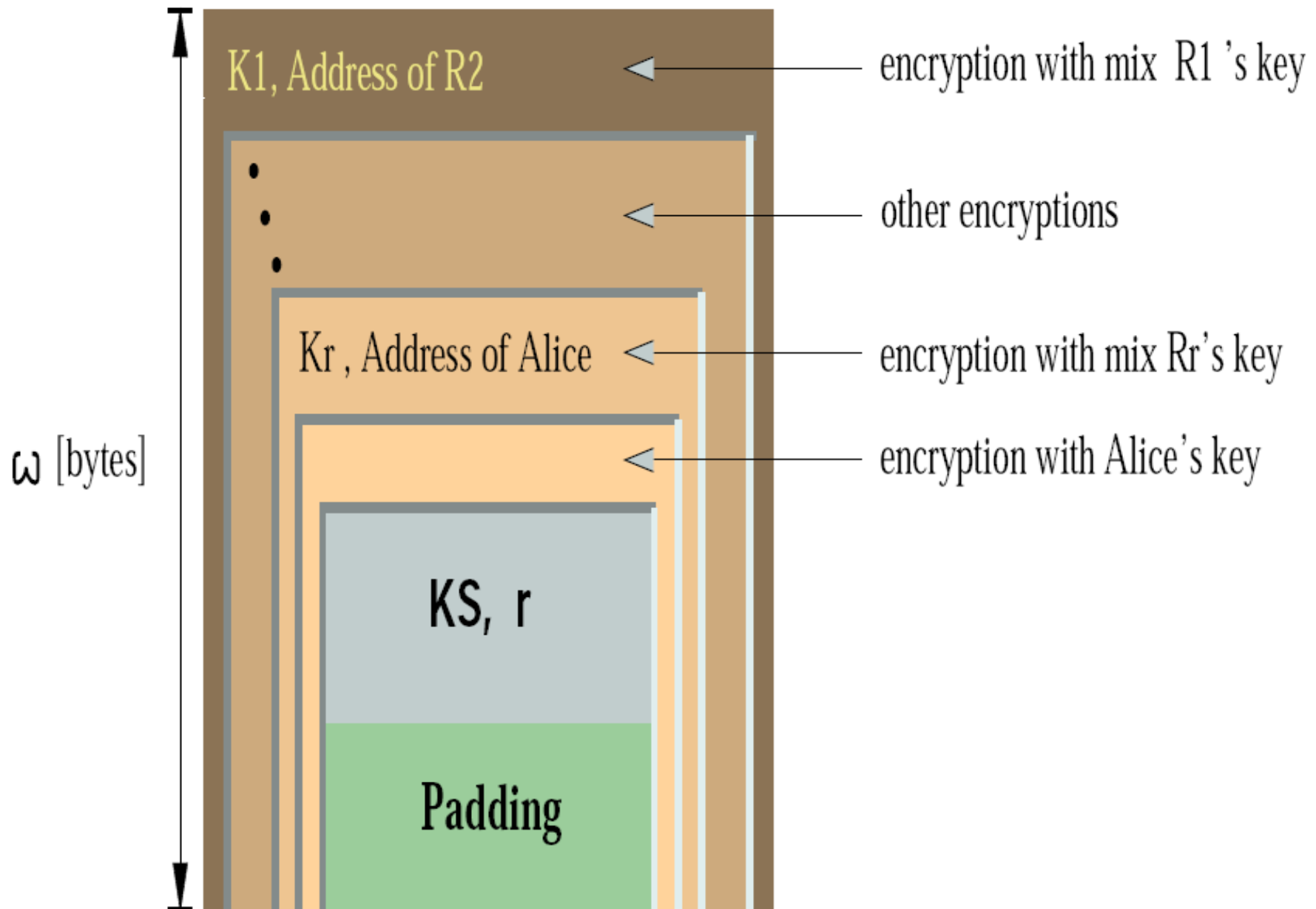
Return path



RPI

$$y_r = \mathcal{A}_{R_1}, E_{R_1}(K_1, \mathcal{A}_{R_2}, E_{R_2}(K_2, \dots \\ \dots E_{R_r}(K_r, \mathcal{A}_{\text{Alice}}, E_{\text{Alice}}(KS, r)) \dots)).$$

Return path information



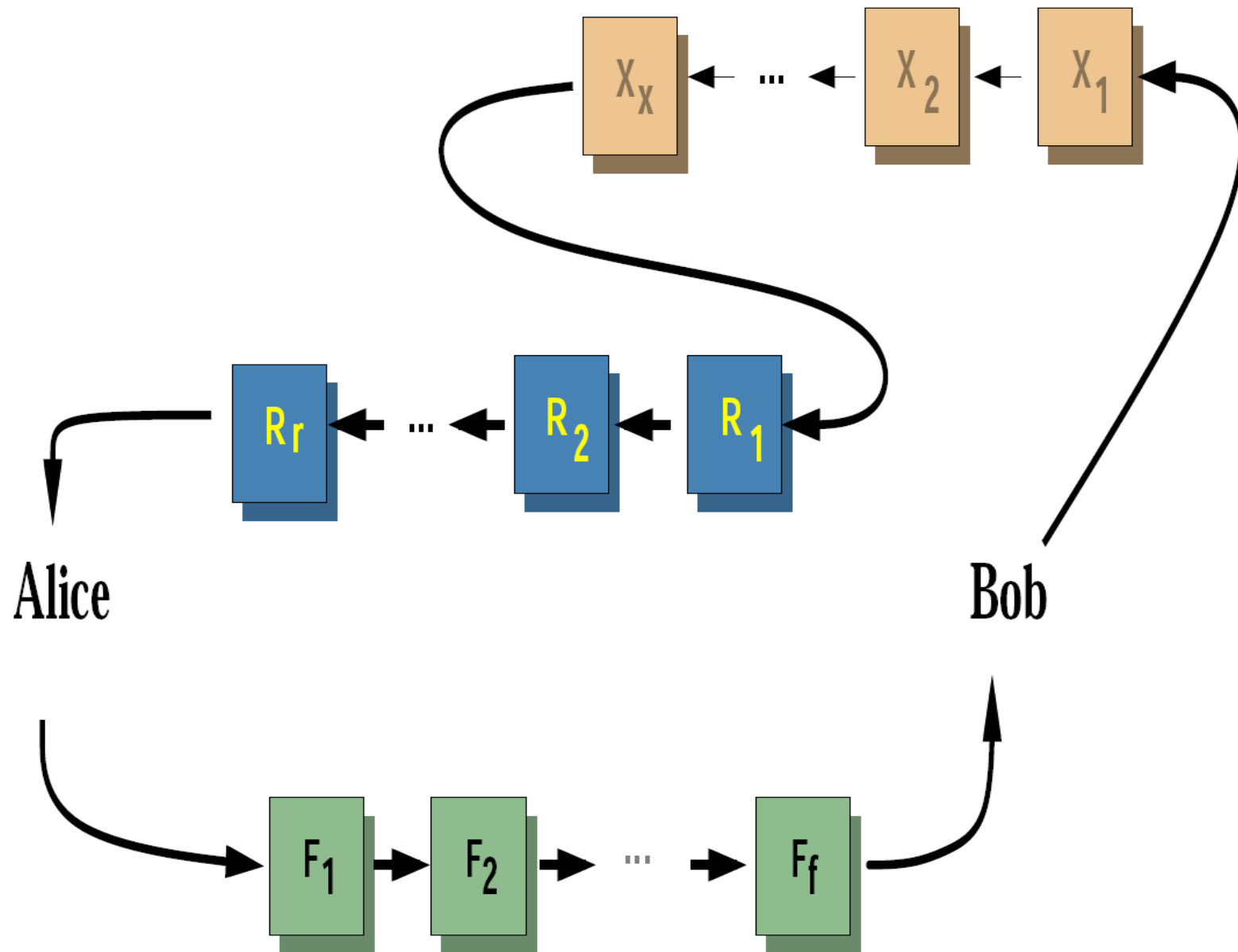
Reply message

Email (SMTP) Header

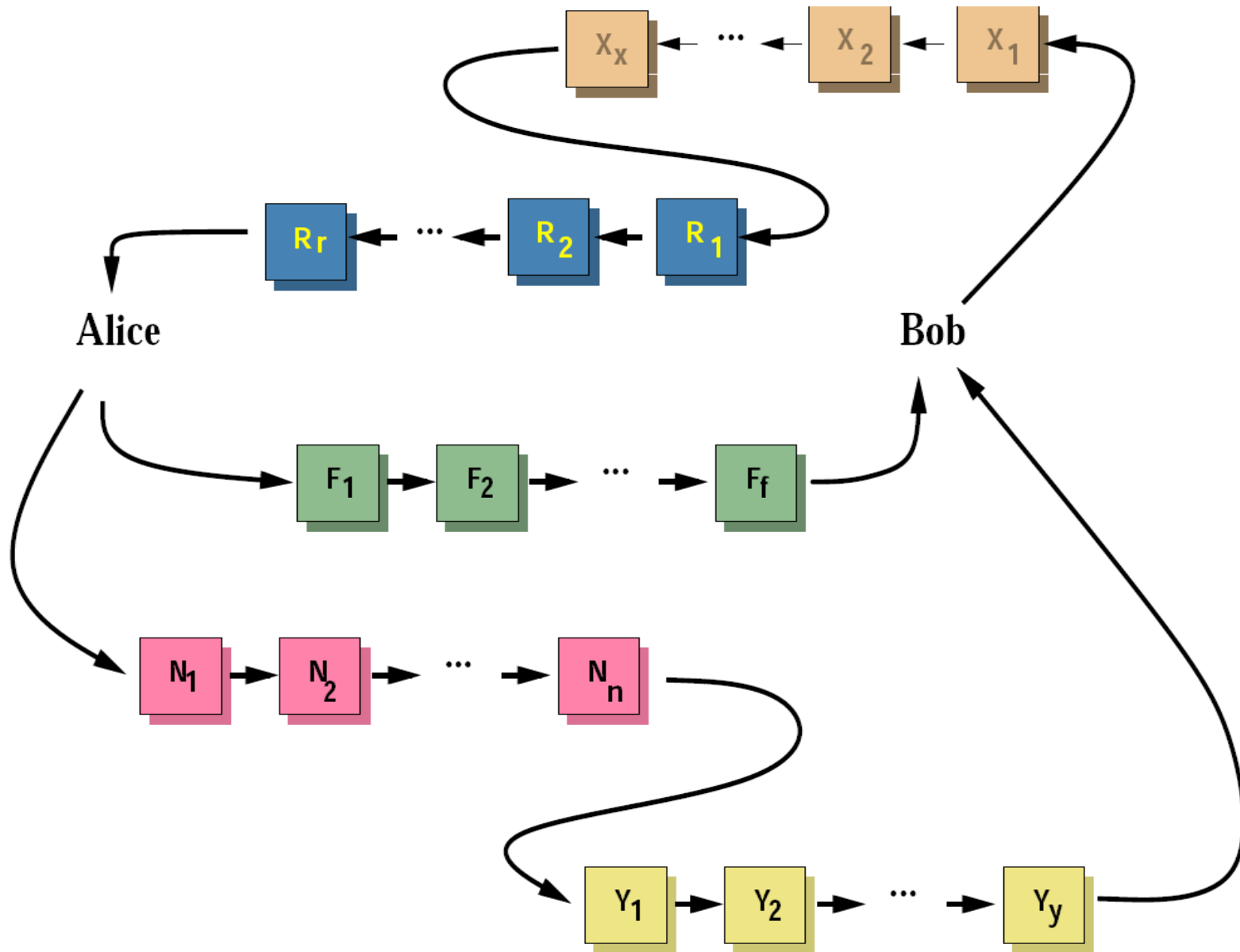
RPI

Message Body

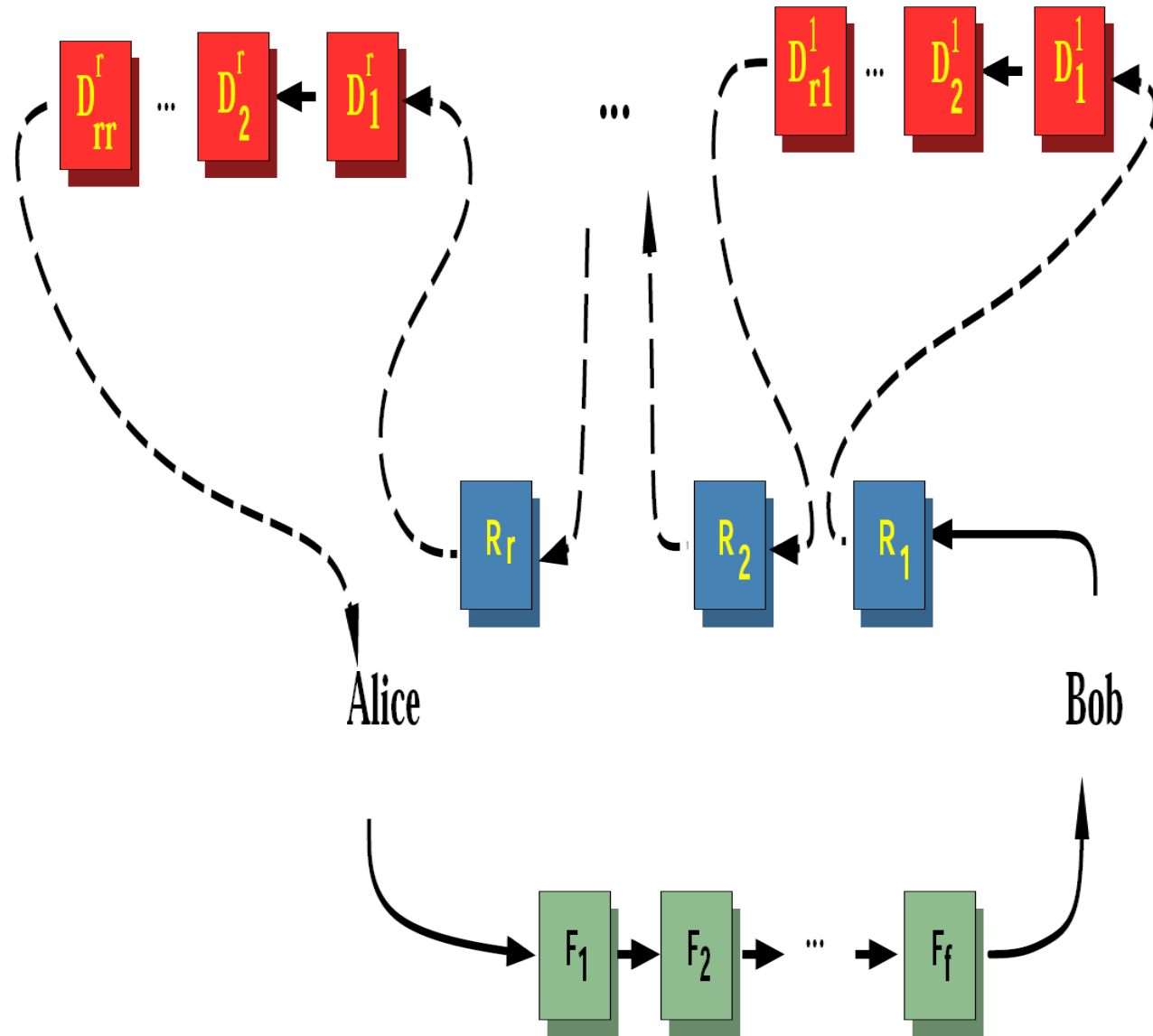
Two-way anonymity



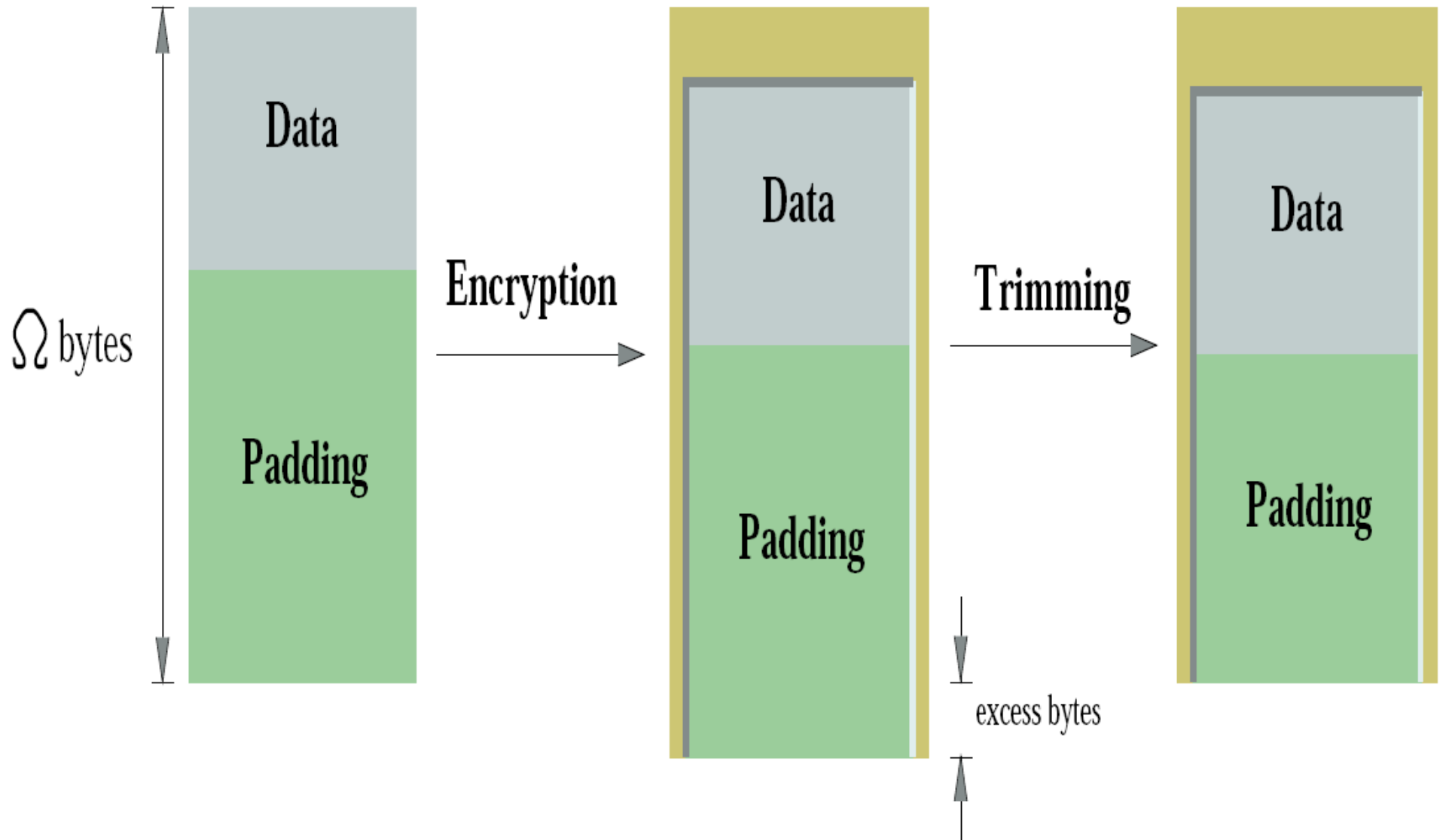
Security of replies



Inter-mix detours



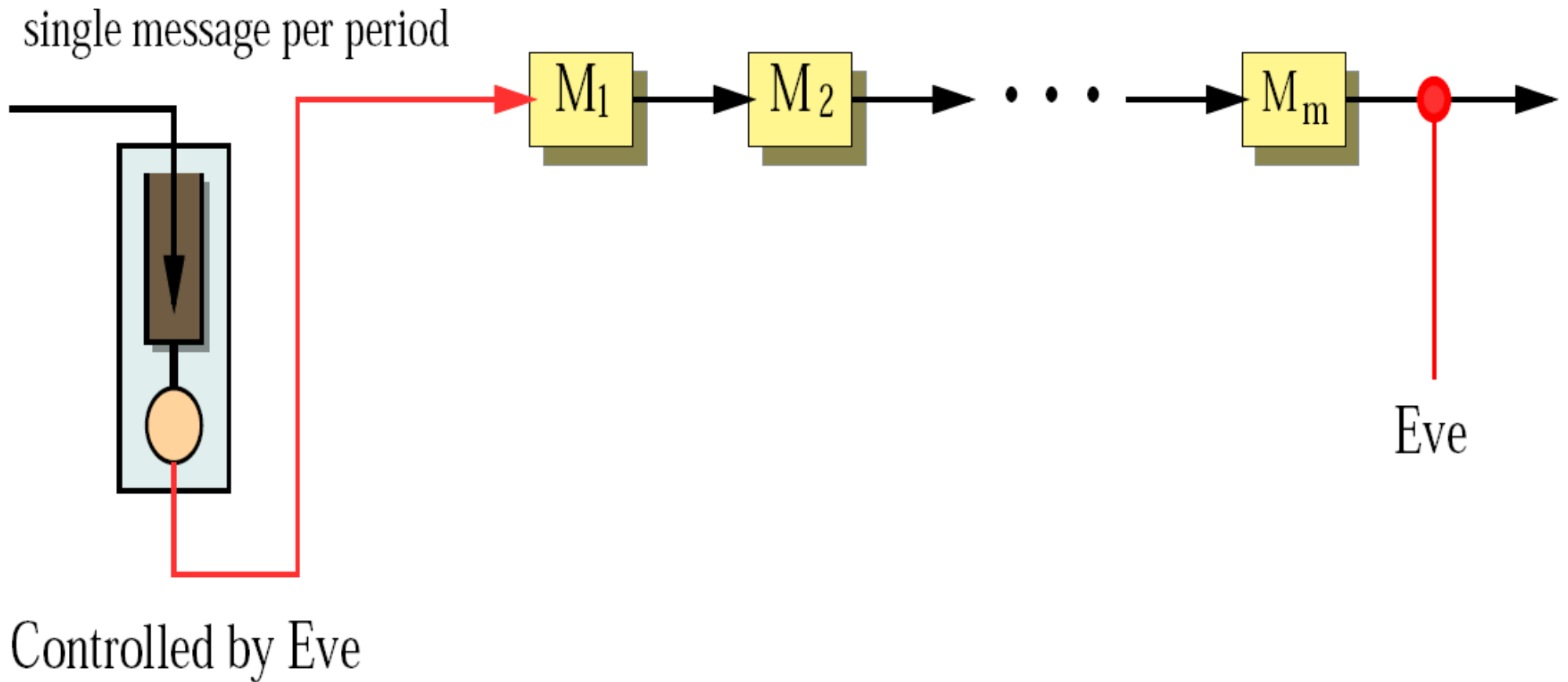
Keeping message size constant



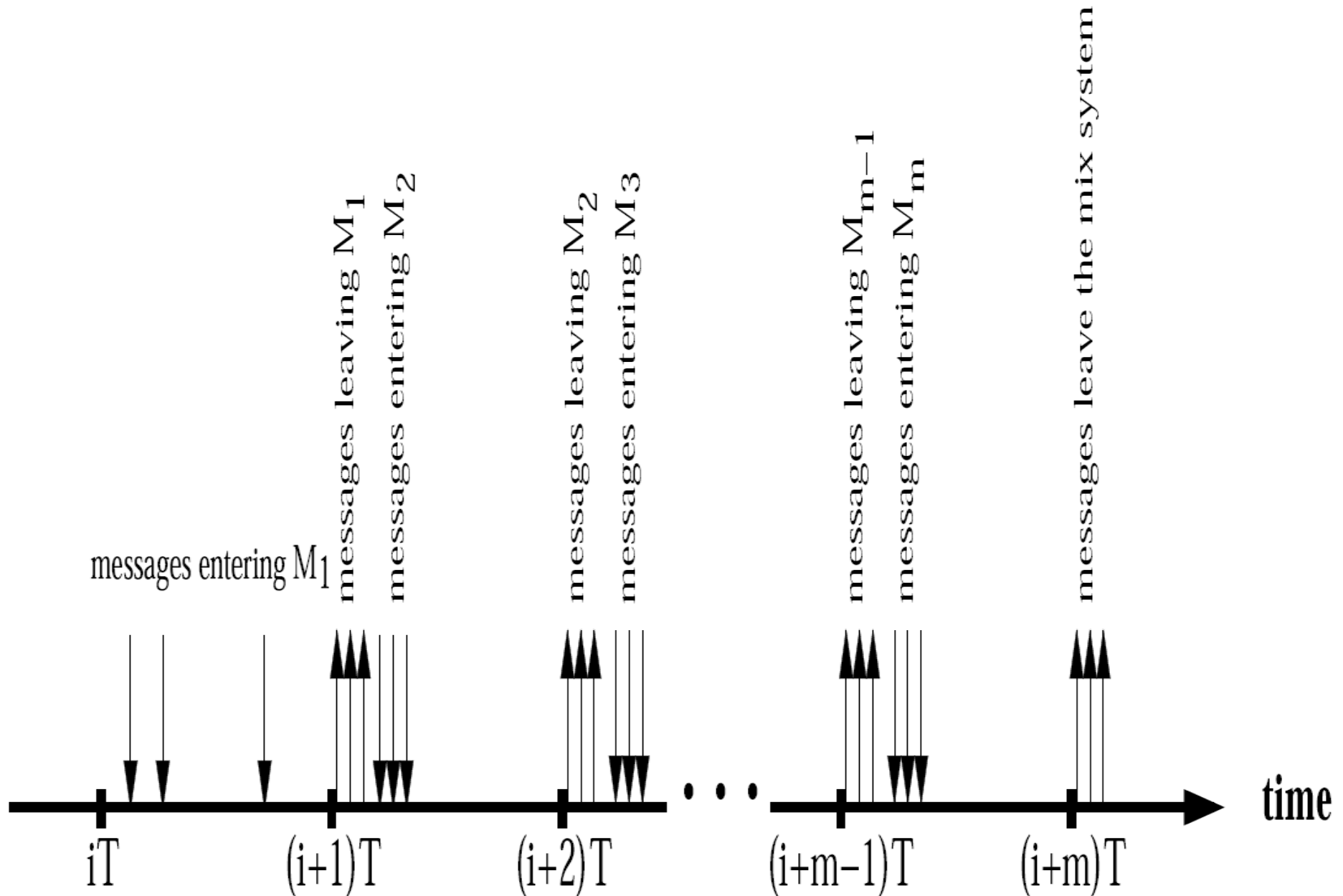
Heeding anonymity

- Fixed-path systems
- System staunchness, miss & guess factors
- Quest for confusion
 - Probabilistic deferment
 - Hybrid approach

Trickle attack



Interval batching



Probabilistic deferment

$$P\{K = k\} = \binom{m}{k} q^{m-k} d^k \quad \text{where } k = 0, \dots, m,$$

Binomial function

